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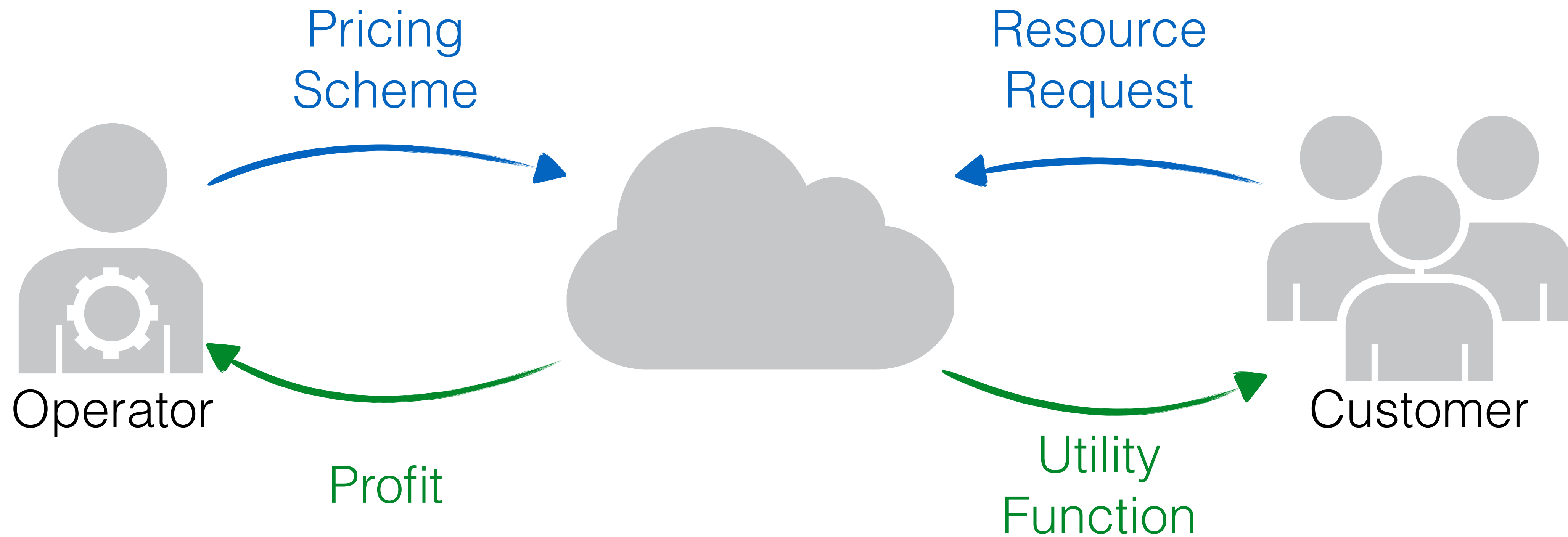
On Pricing Schemes in Data Center Network with Game Theoretic Approach

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Operator and Customer



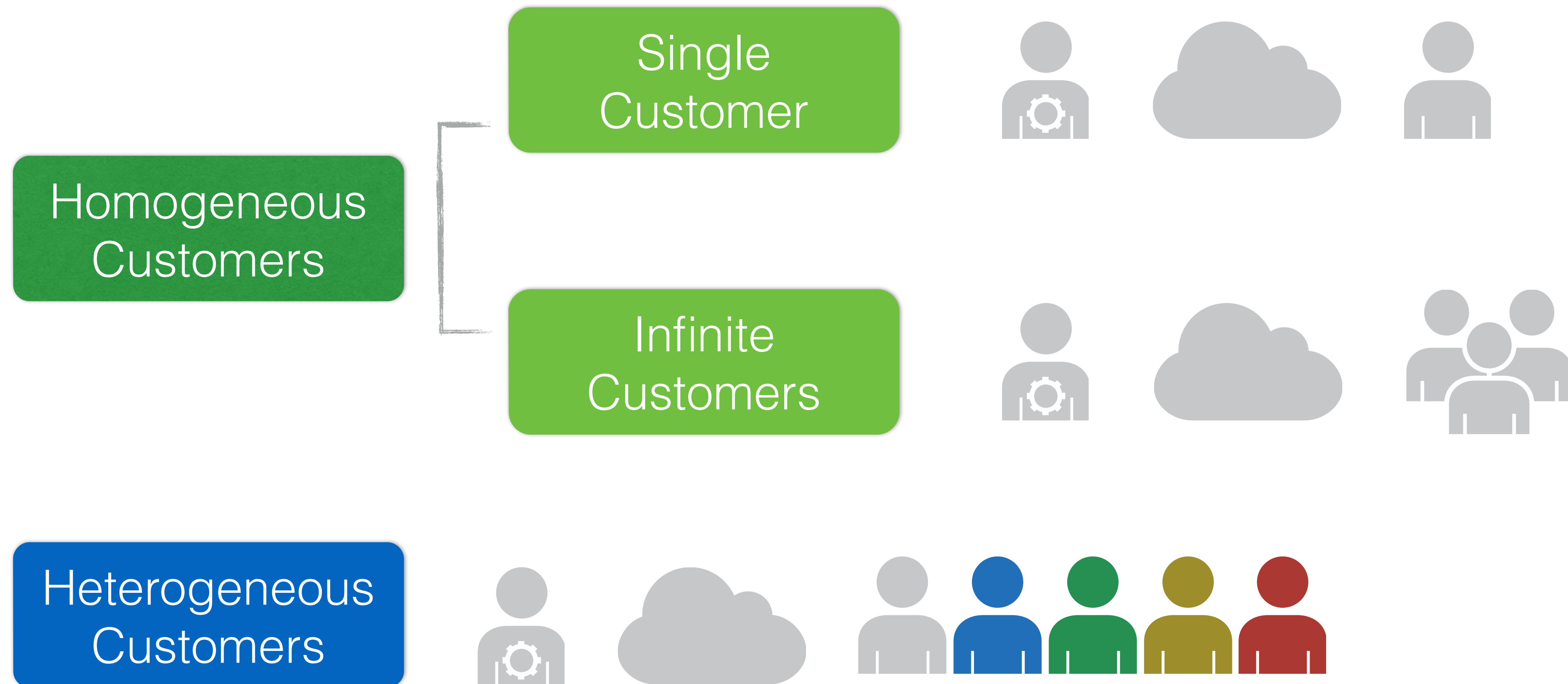
Pricing in DCN

Problem:

- The *operator* announces a price for the resources
- Given this price, the *customer* selfishly determines how many resources it requests

Stackelberg equilibrium will lead to a Pareto-inefficient outcome

Modelling the DCN Operator and Customers



Single Customer

- Customer's utility: $U(d, p) = f(d) - pd$
- Operator's utility: $V(d, p) = pd - vg(d)$

To maximize utility of the customer: $d = f'^{-1}(p) = h(p)$

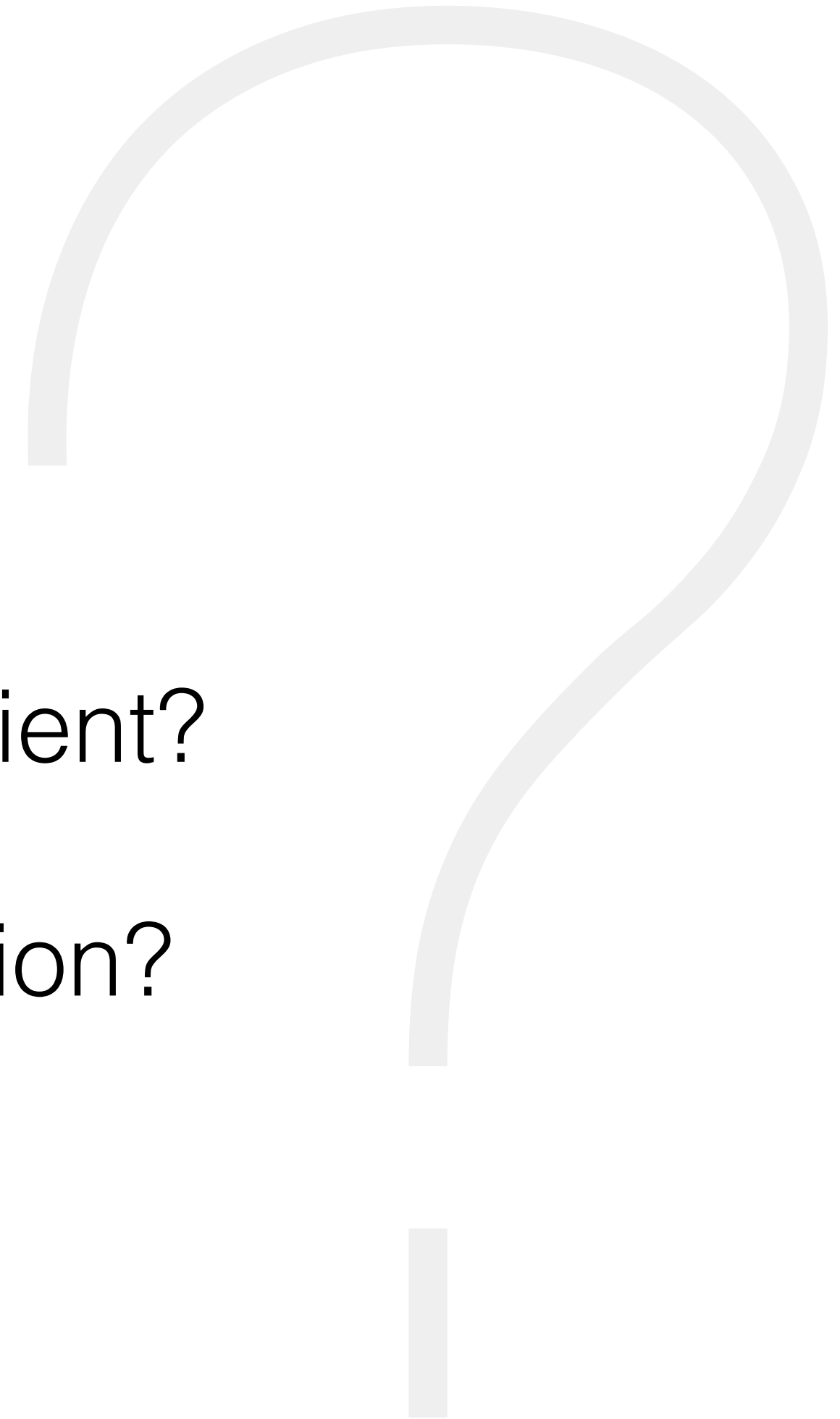
$$V(d(p), p) = ph(p) - vg(h(p))$$

Therefore, operator can maximize its utility by solving:

$$\frac{dV(d(p), p)}{dp} = h(p) + ph'(p) - vg'(h(p))h'(p) = 0$$

Single Customer

- 1) Does Stackelberg equilibrium always exist?
- 2) Is the outcome at Stackelberg equilibrium Pareto-efficient?
- 3) If 2) is not the case, how to get a Pareto-efficient solution?



Single Customer

Example 1

- set $f(d) = \ln d$, $g(d) = d^\alpha$

$$\frac{dV(d(p), p)}{dp} = \frac{1}{p} + p\left(-\frac{1}{p^2}\right) + v\alpha p^{-\alpha-1} = v\alpha p^{-\alpha-1} > 0$$

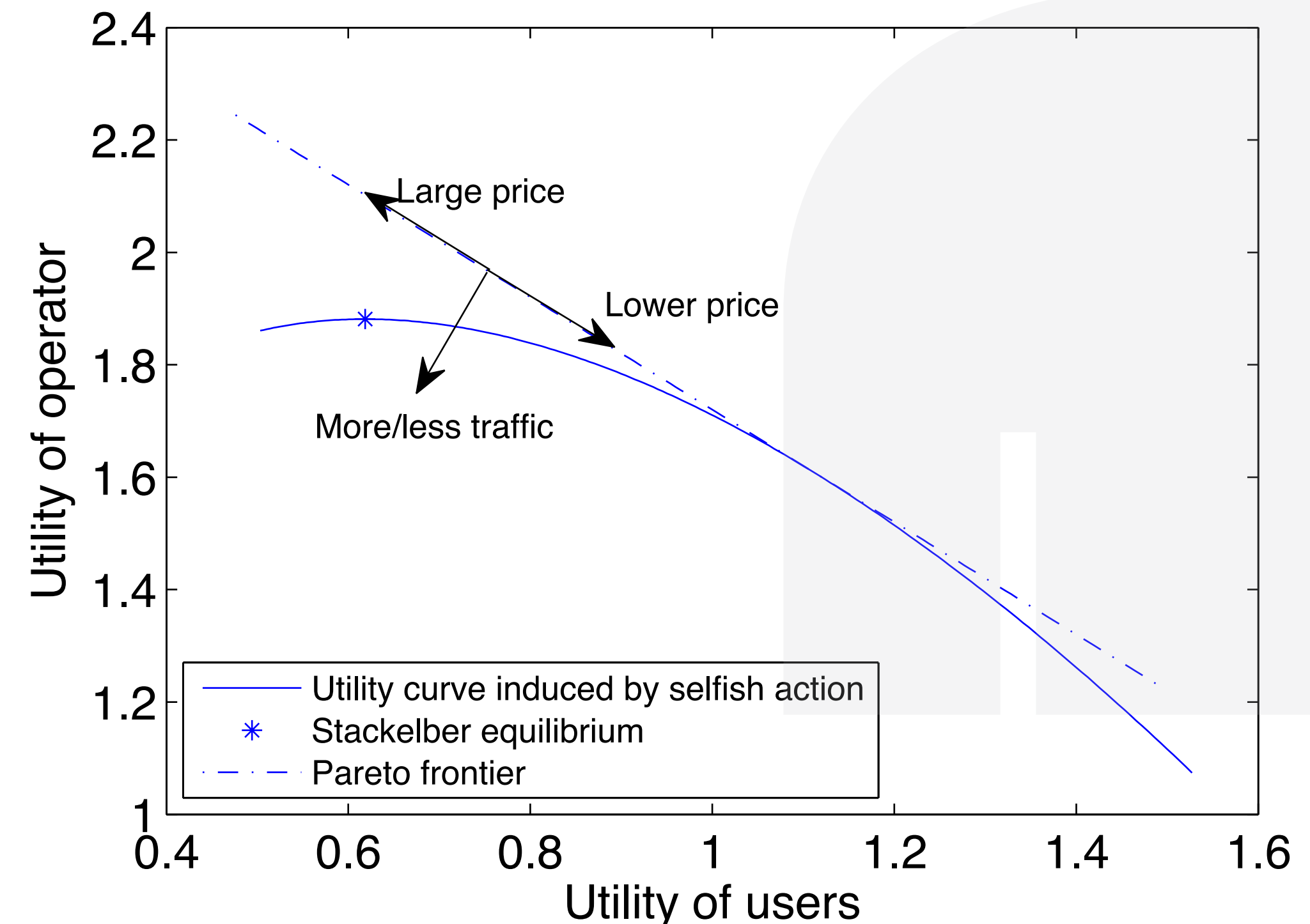
- Stackelberg equilibrium always does not exist

Example 2

- set $f(d) = M - Me^{-d}$

operator will choose a price p such that

$$\ln \frac{p}{M} + 1 = v\alpha \frac{1}{p} \left(-\ln \frac{p}{M}\right)^{\alpha-1}$$



Single Customer

Theorem 1

Assume the demand quantity sent by customer at Stackelberg equilibrium and Pareto efficient solution are $d^{(s)}$ and $d^{(o)}$, respectively, then $d^{(s)} < d^{(o)}$.

Proof:

$$\begin{aligned}
 & U(d^{(o)}, p^{(o)}) - U(d^{(s)}, p^{(s)}) \\
 &= f(d^{(o)}) - f(d^{(s)}) - [f'(d^{(o)})d^{(o)} - f'(d^{(s)})d^{(s)}] \\
 &< f'(\xi)[d^{(o)} - d^{(s)}] - f'(d^{(s)})[d^{(o)} - d^{(s)}] \\
 &= [f'(\xi) - f'(d^{(s)})][d^{(o)} - d^{(s)}] < 0 \quad \dots \dots \dots \textcircled{1}
 \end{aligned}$$

$$U(d^{(o)}, p^{(o)}) + V(d^{(o)}, p^{(o)}) \geq U(d^{(s)}, p^{(s)}) + V(d^{(s)}, p^{(s)}) \quad \dots \textcircled{2}$$

$$V(p^{(o)}, d^{(o)}) \geq V(p^{(s)}, d^{(s)}) \quad \dots \dots \dots \textcircled{3}$$

$$d^{(s)} = h(p^{(s)}) \leq h(p^{(o)}) = d^{(o)} \quad \dots \textcircled{4}$$

Single Customer

The price for d units of resource: $B(d, p) = P + pd \dots\dots\dots \mathbf{1}$

Theorem 2 Assume $T(d) = U(d, p) + V(d, p)$

$$\begin{cases} p = p^{(o)} \\ P = V(d^{(s)}, p^{(s)}) - V(d^{(o)}, p^{(o)}) + \frac{1}{2}(T(d^{(o)}) - T(d^{(s)})) \end{cases} \quad \begin{array}{l} \text{Pareto efficient and} \\ \text{Stackelberg equilibrium} \end{array}$$

Proof: From $\mathbf{1}$, $U_a(d, p) = f(d) - (P + p^{(o)}d)$

$$\begin{aligned} U_a(d^{(o)}, p^{(o)}) &= f(d^{(o)}) - (P + p^{(o)}d^{(o)}) = U(d^{(o)}, p^{(o)}) - [V(d^{(s)}, p^{(s)}) - V(d^{(o)}, p^{(o)}) + \frac{1}{2}(T(d^{(o)}) - T(d^{(s)}))] \\ &= \frac{1}{2}(T(d^{(o)}) + T(d^{(s)})) - V(d^{(s)}, p^{(s)}) > T(d^{(s)}) - V(d^{(s)}, p^{(s)}) = U(d^{(s)}, p^{(s)}) \end{aligned}$$

$$\begin{aligned} V_a(d^{(o)}, p^{(o)}) &= (P + p^{(o)}d^{(o)}) - vg(d^{(o)}) = P + V(d^{(o)}, p^{(o)}) \\ &= V(d^{(s)}, p^{(s)}) + \frac{1}{2}(T(d^{(o)}) - T(d^{(s)})) > V(d^{(s)}, p^{(s)}) \end{aligned}$$

Infinite Customers

Lemma 1

- Customer's utility: $U(d, p) = f(d) / d - p$
- Operator's utility: $V(d, p) = pd - vg(d)$

When the operator announces a larger price, there will be less demand sent into the network.

Proof:

$$\left. \begin{array}{l} f(d) = pd \\ f'(d)d' = d + pd' \end{array} \right\} \longrightarrow d' = \frac{d^2}{f'(d)d - f(d)}$$

Since $f(0) = 0$,

$$f'(d)d - f(d) = f'(d)d - f(d) + f(0) = [f'(d) - f'(d^*)]d \quad (d^* \in (0, d))$$

Thus, $d' < 0$ and d is a decreasing function of p

Infinite Customers

Lemma 2

The demand quantity d is a concave function of price of each unit resource, if and only if $f''(d)d' > 2$.

Proof:

$$f'(d)d' = d + pd'$$



derivation

$$f''(d)(d')^2 + f'(d)d'' = 2d' + pd''$$

$$d'' = \frac{2d' - f''(d)(d')^2}{f'(d) - p}$$

$$f'(d) - p = d / d'$$



$$\begin{aligned} d'' &= [2(d')^2 - f''(d)(d')^3] / d \\ &= [2 - f''(d)d'](d')^2 / d \end{aligned}$$

Infinite Customers

Lemma 3

$f''(d)d' > 2$ $V(d(p), p)$ is a concave function of p .

Proof:

$f(d)$ is a concave function

$g(d)$ is a convex function

$V(d, p) = f(d) - vg(d)$ is a concave function of d

d is a decreasing and concave function of p .

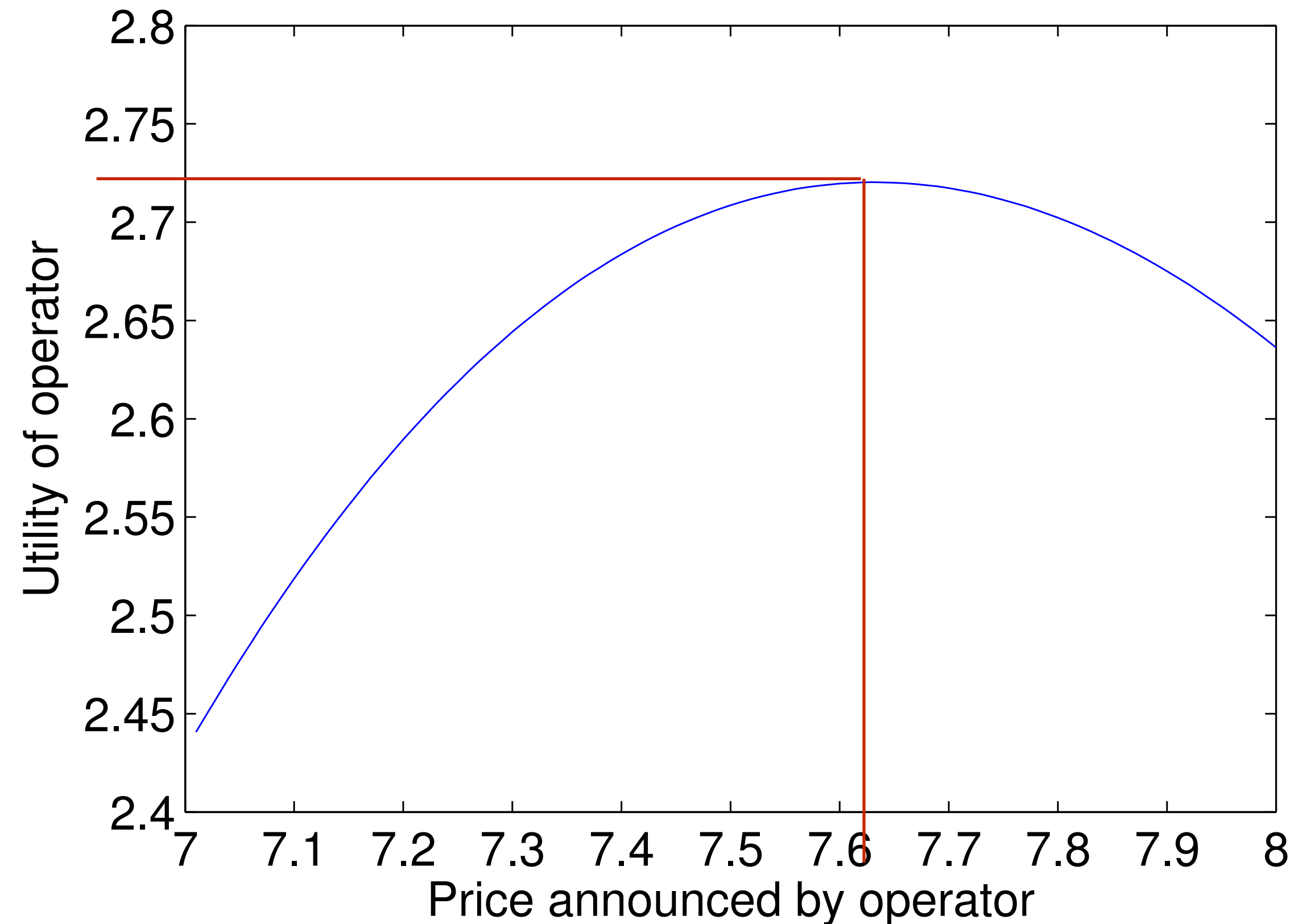


Infinite Customers

Example 3

Utility of customers from one unit of resource:

- set $f(d) = M - Me^{-d}$



search such Stackelberg equilibrium if $V(d(p), p)$ is a concave function of p .

Heterogeneous Customer

- QoS requirement of a customer s
- total demand quantity in the network is D
- distribution density function $q(s)$
- Customer's utility: $U(s, p) = D \int_s^{u^{-1}(p)} u(x)q(x) dx - pD \int_s^{u^{-1}(p)} q(x) dx$
- Operator's utility: $V(s, p) = p \int_s^{u^{-1}(p)} Dq(x) dx - vg(\int_s^{u^{-1}(p)} Dq(x) dx)$
- the QoS requirement constraint $D \int_s^{u^{-1}(p)} q(x) dx = sC$
- marginal utility $u(s)$
- network resource capacity is C

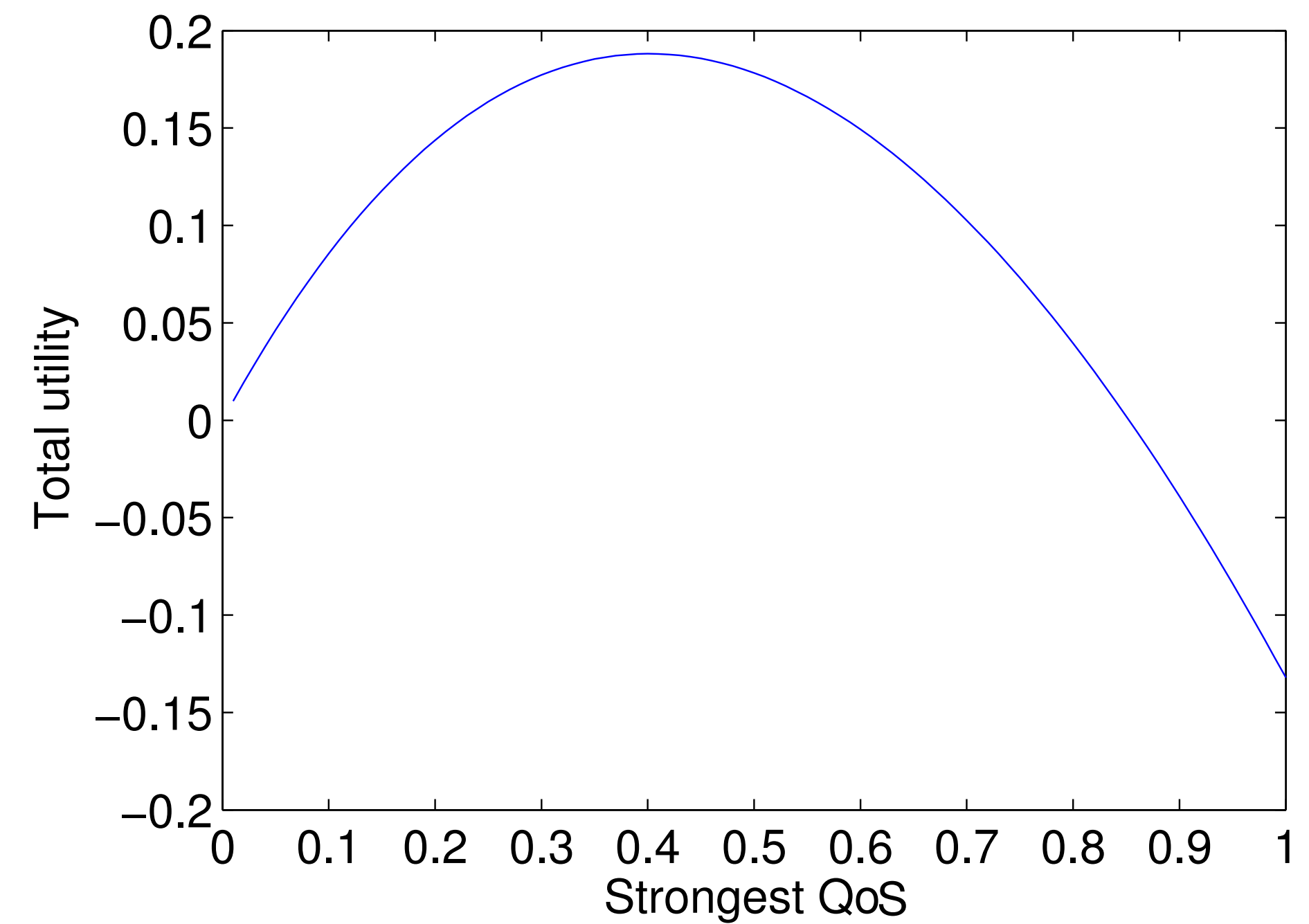
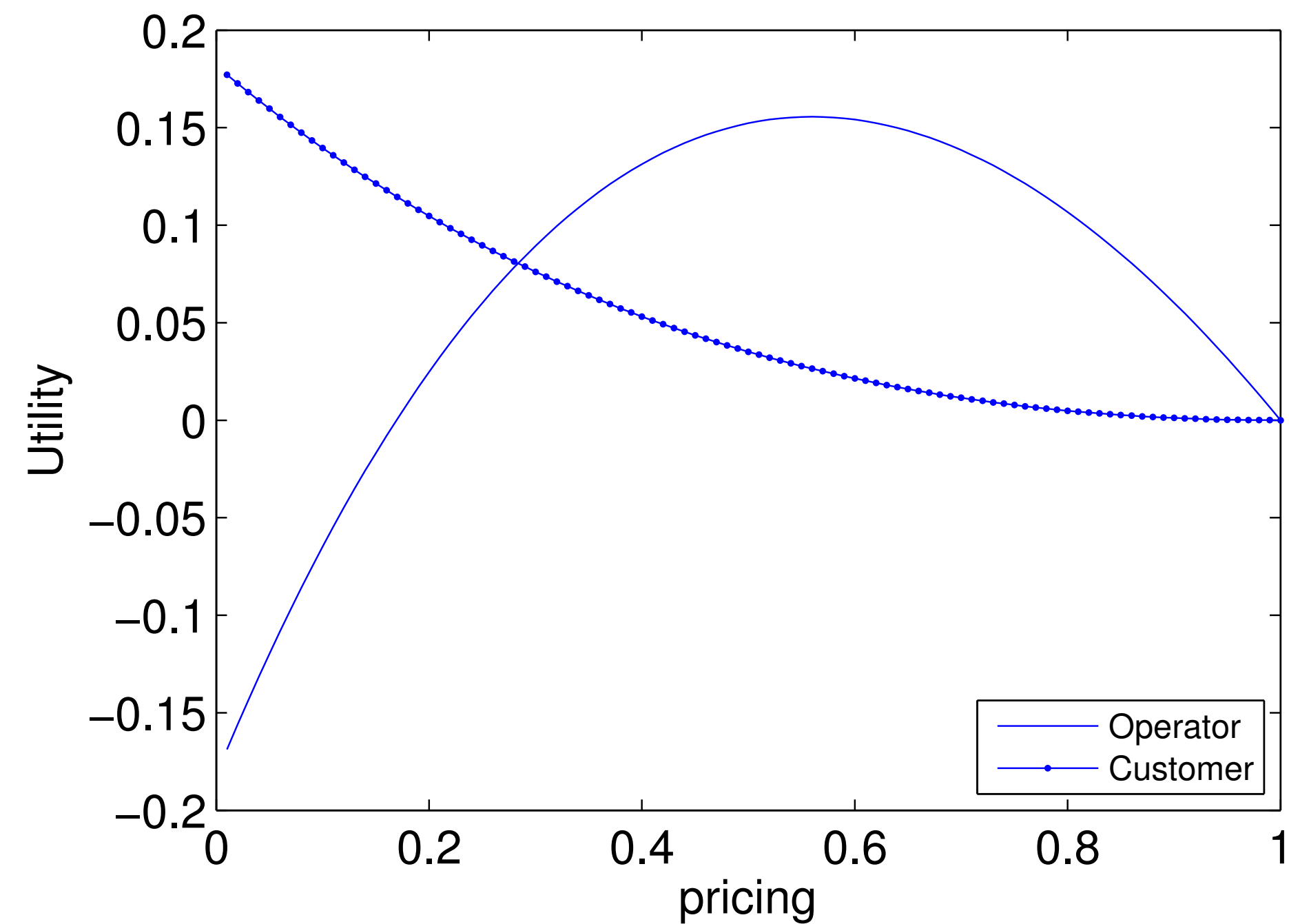
$$V(h(p), p) = p \int_{h(p)}^{u^{-1}(p)} Dq(x) dx - vg(\int_{h(p)}^{u^{-1}(p)} Dq(x) dx)$$

$$pD[I'(p)q(I(p)) - h'(p)q(h(p))] + D \int_{h(p)}^{I(p)} q(x) dx =$$

$$vDg'(D \int_{h(p)}^{I(p)} q(x) dx)[I'(p)q(I(p)) - q(h(p))h'(p)]$$

Heterogeneous Customer

Example 4 set $u(s) = e^{-s}$, $q(s) = e^{-s}$, $g(d) = d^2$, $D = 2$, $C = 1$, $v = 1$



Heterogeneous Customer

To enable Paris Metro Pricing (PMP):

- How to determine the price of each subnetwork?
- How to assign resource to each subnetworks?

Heterogeneous Customer

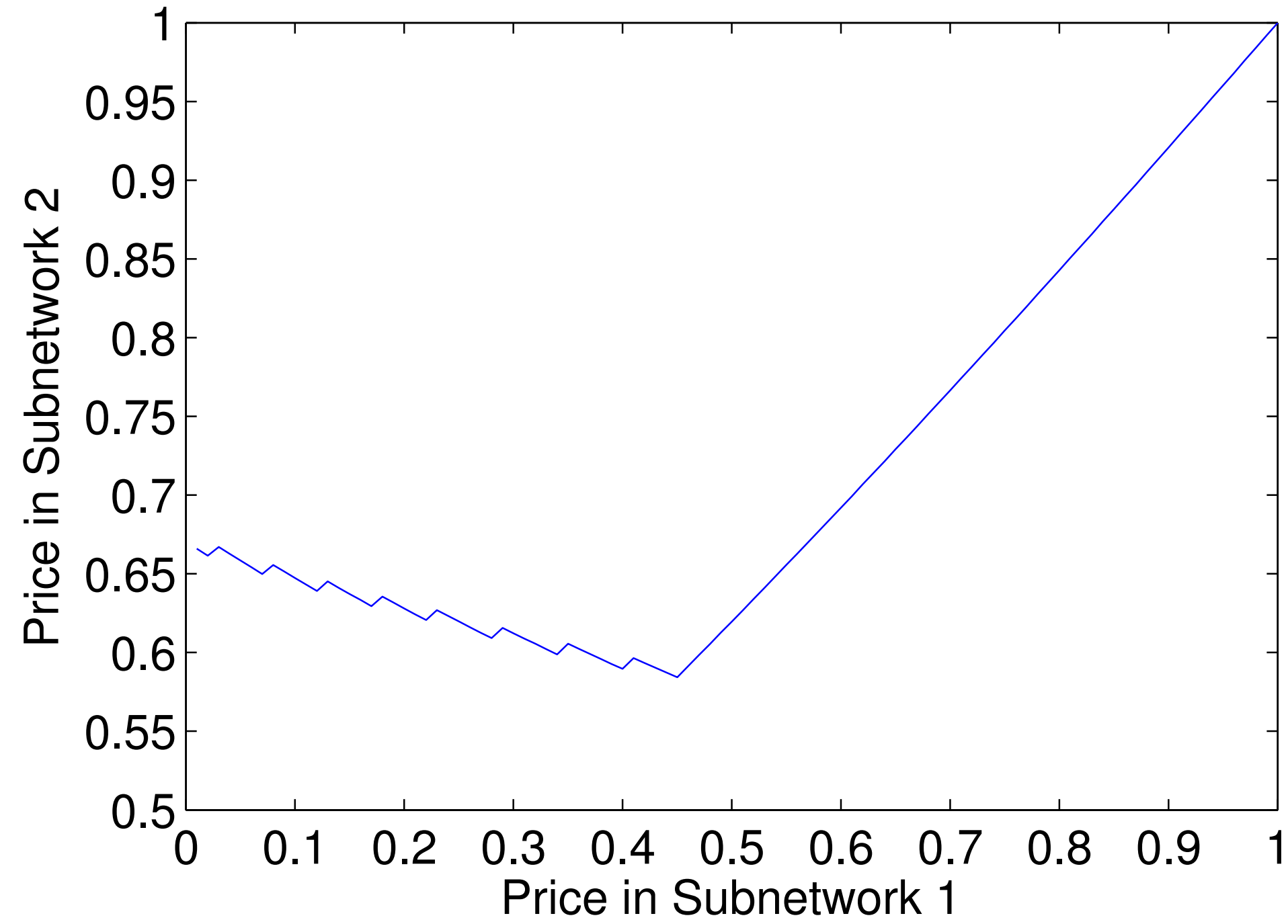
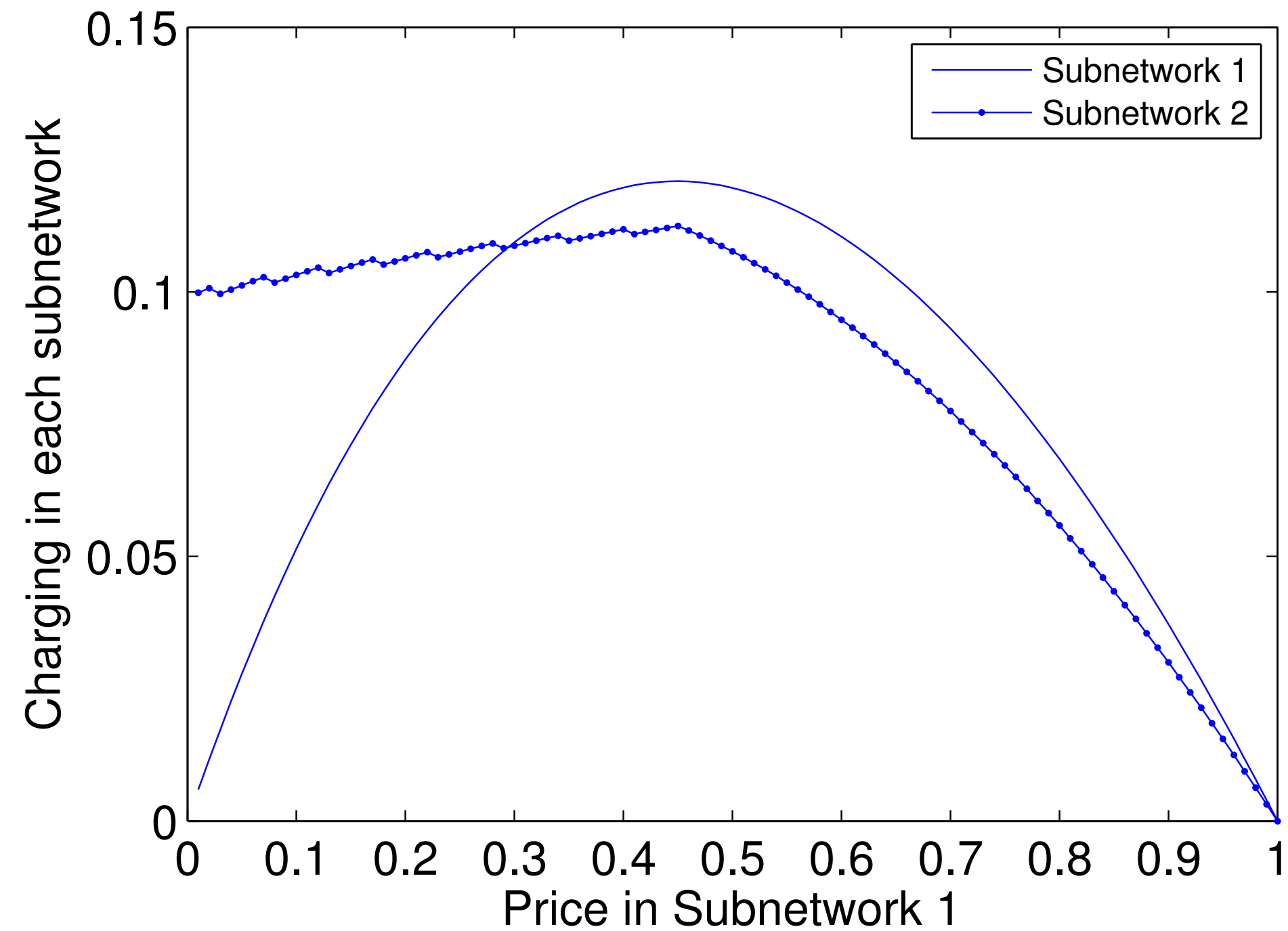
- Leveraging Paris Metro Pricing (PMP)

- Customer's utility:
$$U(s_1, s_2, p_1, p_2) = D \int_{s_1}^{u^{-1}(p_1)} u(x)g(x) dx + D \int_{s_2}^{\min(s_1, u^{-1}(p_2))} u(x)g(x) dx - p_1 D \int_{s_1}^{u^{-1}(p_1)} q(x) dx - p_2 D \int_{s_2}^{\min(s_1, u^{-1}(p_2))} q(x) dx$$

- Operator's utility:
$$V(s_1, s_2, p_1, p_2) = p_1 D \int_{s_1}^{u^{-1}(p_1)} q(x) dx + p_2 D \int_{s_2}^{\min(s_1, u^{-1}(p_2))} q(x) dx - vg(D \int_{s_1}^{u^{-1}(p_1)} q(x) dx + D \int_{s_2}^{\min(s_1, u^{-1}(p_2))} q(x) dx)$$

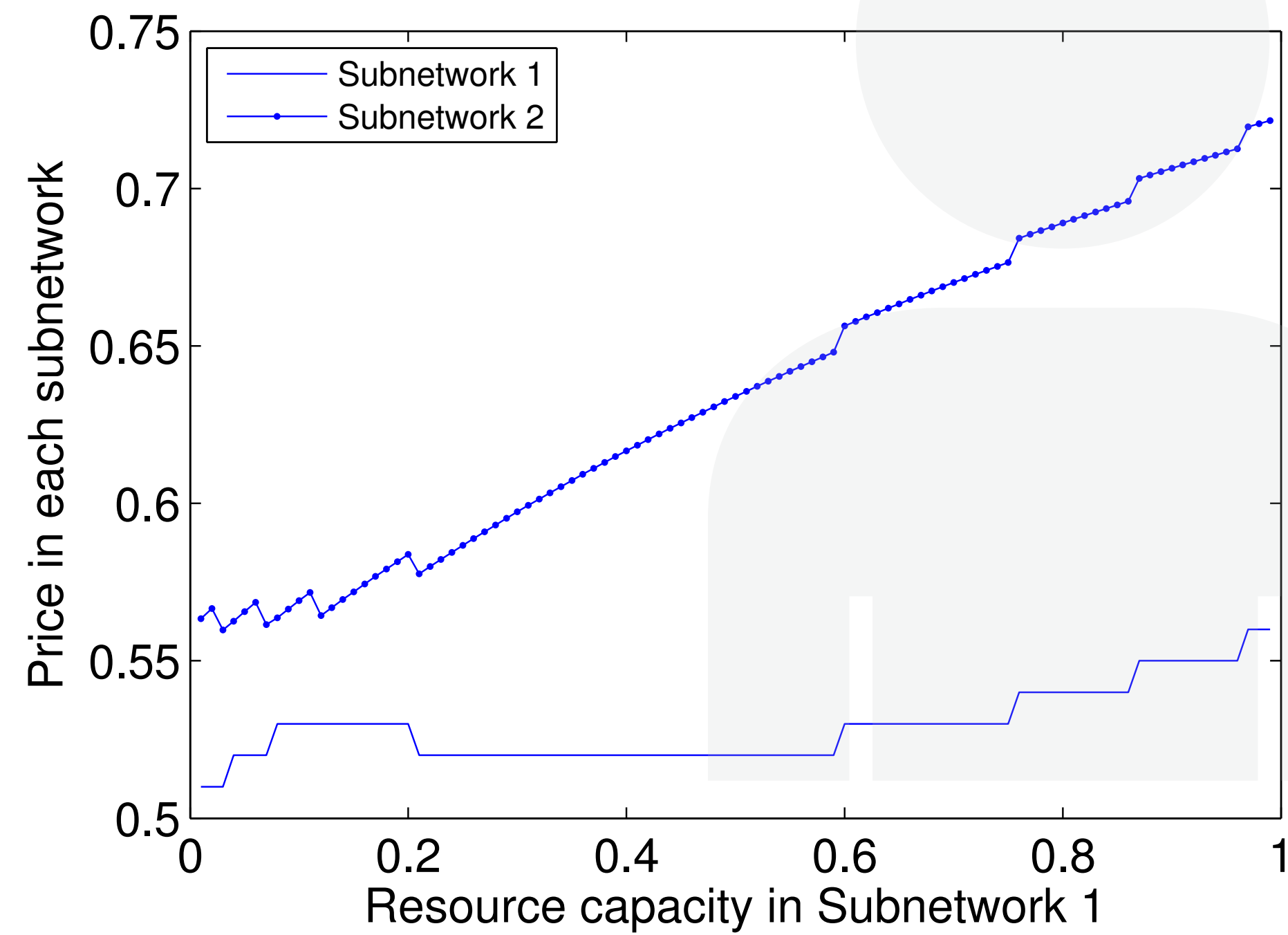
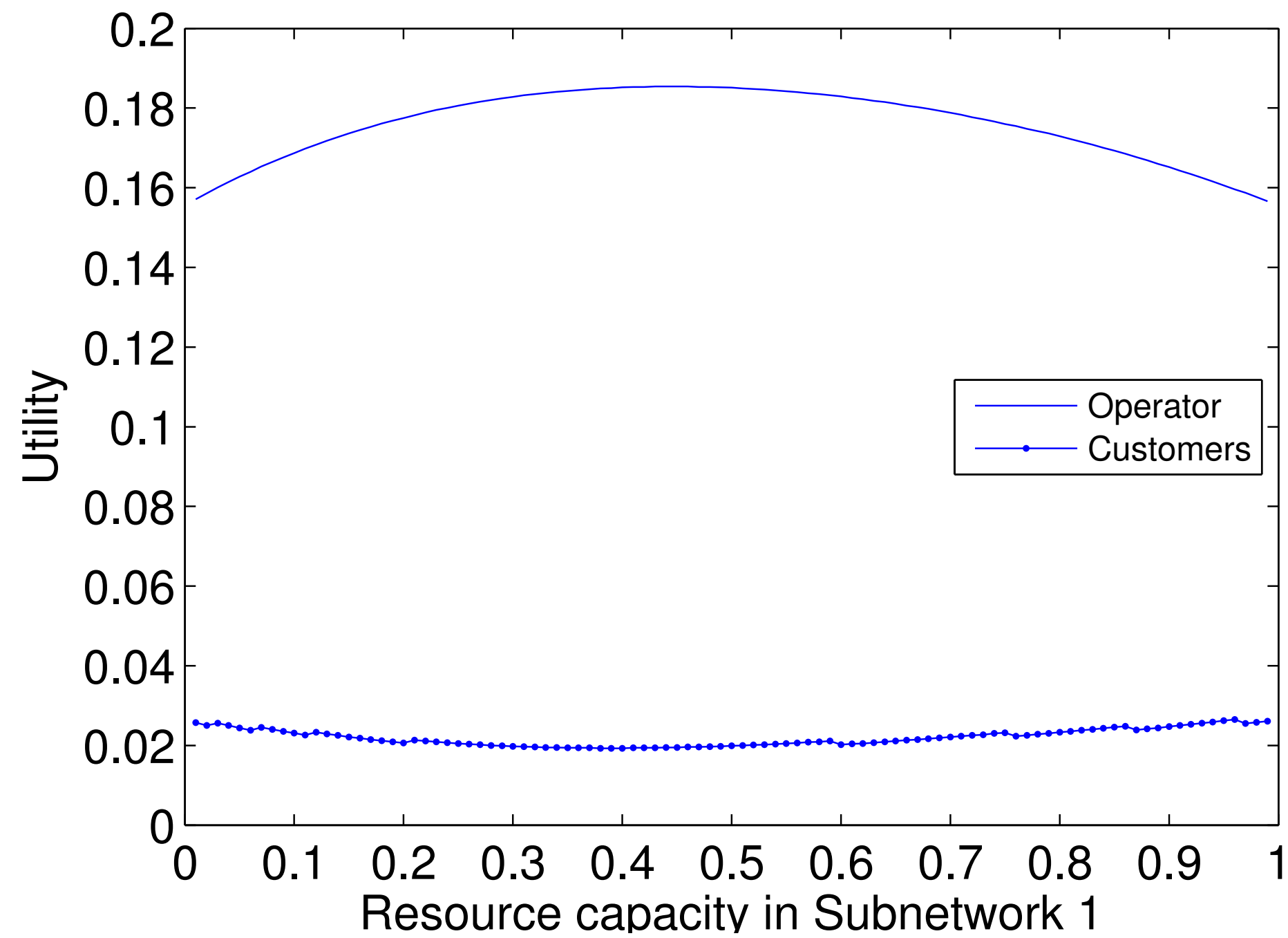
Heterogeneous Customer

Example 5



Heterogeneous Customer

Example 6



Conclusion

- Modelling the pricing problem in DCN as a Stackelberg game.
- Classify the market based on customers.
- For homogeneous and heterogeneous customers, Pareto-efficient solution at Stackelberg equilibrium.
- Introduce PMP scheme to heterogeneous customers case.



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Thank You

